

Lecture six

**Topics that must be covered in this lecture:**

- **Finite automata with output**
- **Moore Machine**
- **Mealy Machine**
- **Equivalence between Moore and Mealy Machines**

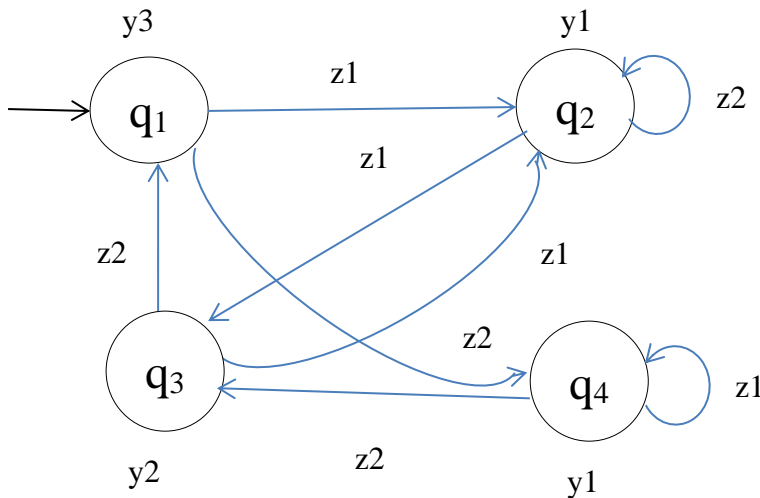
**Finite automata with output**

One limitation of the finite automata as we have defined it is that its output limited to a binary signal: " accept"/"reject". Models in which the output chosen from other alphabet have been considered. There are two distinct approaches; the output may be associated with the state (called a Moore machine) or with the transition ( called a Mealy machine).

**Moore Machine**

A Moore machine is a six- tuple  $(Q, \Sigma, \Delta, \delta, \lambda, q_0 = \text{start state})$ , where  $Q, \Sigma, \delta, q_0$  are as in the DFA.  $\Delta$  is the output alphabet and  $\lambda$  is the mapping from  $Q$  to  $\Delta$  giving the output associated with each state. The output of  $M$  in response to input  $a_1 a_2 \dots a_n, n \geq 0$ , is  $\lambda(q_0) \lambda(q_1) \dots \lambda(q_n)$ , where  $q_0, q_1, \dots, q_n$  is the sequence of states such that  $\delta(q_{i-1}, a_i) = q_i$  for  $1 \leq i \leq n$ . Note that any Moore machine gives output  $\lambda(q_0)$  in response to input  $\epsilon$ .

**Example:** Give the formal description for the following Moore machine  $M_1$  pictured below:



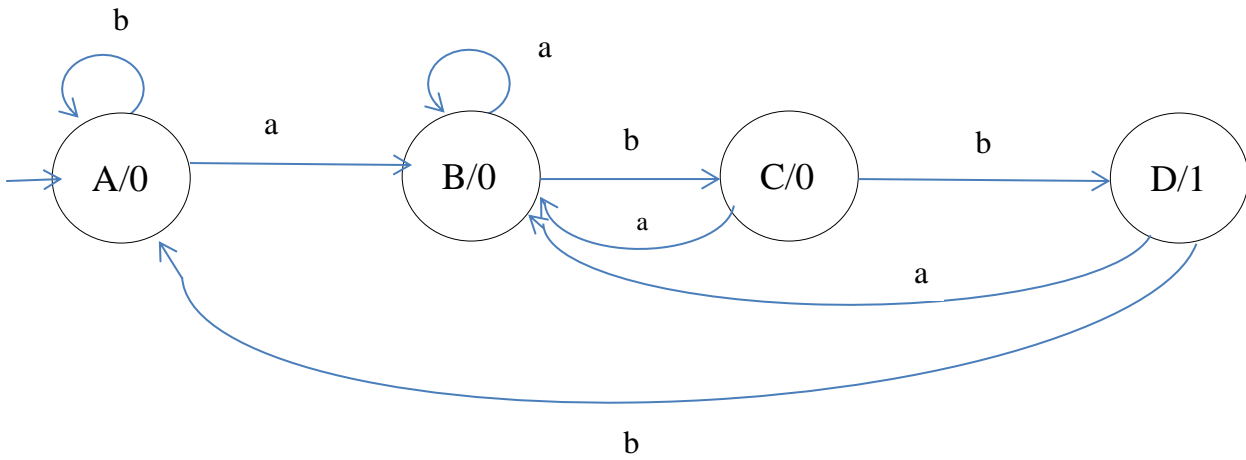
**Sol:**

$$M_1 = (\{q_1, q_2, q_3, q_4\}, \{z_1, z_2\}, \{y_1, y_2, y_3\}, \delta, \lambda, q_1)$$

$$\delta: Q \times \Sigma \rightarrow Q, \lambda: Q \rightarrow \Delta$$

output state	y3 q1	y1 q2	y2 q3	y1 q4
input				
z1	q2	q3	q2	q4
z2	q4	q2	q1	q3

**Example2:** Give the formal description for the following Moore machine M2 pictured below, and give the output to the input string abbabb

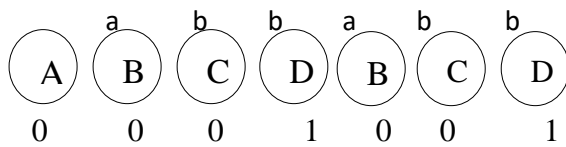


Sol:

-  $M2 = (\{A, B, C, D\}, \{a, b\}, \{0, 1\}, \delta, \lambda, A)$

output state input	0 A	0 B	0 C	1 D
a	B	B	B	B
b	A	C	D	A

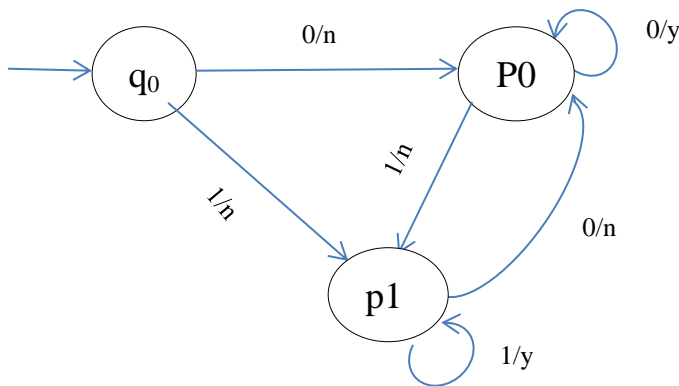
- OUTPUT:



**Mealy Machine**

A mealy machine is also a six – tuple  $(Q, \Sigma, \Delta, \delta, \lambda, q_0=\text{start state})$ , where all is as in the Moore machine, except that  $\lambda$  maps  $Q \times \Sigma \rightarrow \Delta$ . That is,  $\lambda(q,a)$  gives the output associated with the transition from state  $q$  on input  $a$ . the output of  $M$  in response to input  $a_1a_2\dots a_n$  is  $\lambda(q_0,a_1), \lambda(q_1,a_2), \dots, \lambda(q_{n-1},a_n)$ , where  $q_0, q_1, \dots, q_n$  is the sequence of states such that  $\delta(q_{i-1}, a_i) = q_i$  for  $1 \leq i \leq n$ . Note that the sequence has length  $n$  rather than length  $n+1$  as for the Moore machines and on input  $\epsilon$  a Mealy machine gives output  $\epsilon$ .

**Example:** Give the formal description for the following Mealy machine M3 pictured below:



Sol:

$M3 = (\{q_0, p_0, p_1\}, \{0, 1\}, \{n, y\}, \delta, \lambda, q_0)$

$\delta$  – table

Q	q0	p0	p1
$\Sigma$			
0	p0	p0	p0
1	p1	p1	p1

$\lambda$  – table

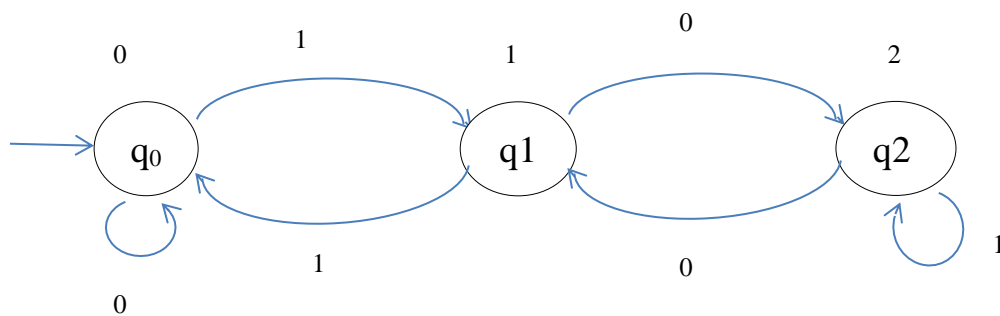
Q	q0	p0
$\Sigma$		
0	n	y
1	n	n

**Equivalence between Moore and Mealy Machines**

**Theorem:** if  $M1 = (Q, \Sigma, \Delta, \delta, \lambda, q_0)$  is a Moore Machine, then there is a Mealy Machine  $M2$  equivalent to  $M1$ .

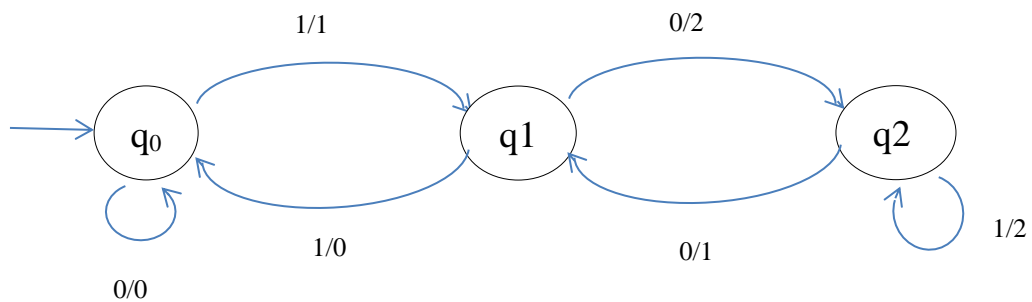
**Proof:** Let  $M2 = (Q, \Sigma, \Delta, \delta, \lambda', q_0)$  and define  $\lambda'(q,a)$  to be  $\lambda(\delta(q,a))$  for all states  $q$  and input symbols  $a$ . then  $M1$  and  $M2$  enter the same sequence of states on the same input, and with each transition  $M2$  emits the output that  $M1$  associates with the state entered.

**Example1:** Construct an equivalence a Mealy Machine from a Moore Machine pictured below.



**SOLUTION:**

A mealy machine equivalence to a Moore machine :



**Method of conversion :**

$$\lambda'(q,a) = \lambda(\delta(q,a))$$

$$\lambda'(q_0,0) = \lambda(\delta(q_0,0)) = \lambda(q_0) = 0$$

$$\lambda'(q_0,1) = \lambda(\delta(q_0,1)) = \lambda(q_1) = 1$$

$$\lambda'(q_1,0) = \lambda(\delta(q_1,0)) = \lambda(q_2) = 2$$

$$\lambda'(q_1,1) = \lambda(\delta(q_1,1)) = \lambda(q_0) = 0$$

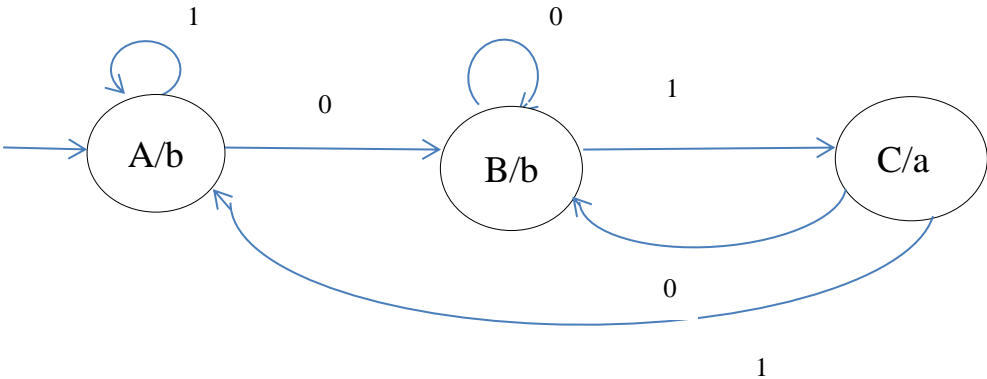
$$\lambda'(q_2,0) = \lambda(\delta(q_2,0)) = \lambda(q_1) = 1$$

$$\lambda'(q_2,1) = \lambda(\delta(q_2,1)) = \lambda(q_2) = 2$$

**Transition table of a mealy machine:**

$\Sigma$ Q	<b>0</b>	<b>1</b>
<b>q0</b>	q0,0	q1,1
<b>q1</b>	q2,2	q0,0
<b>q2</b>	q1,1	q2,2

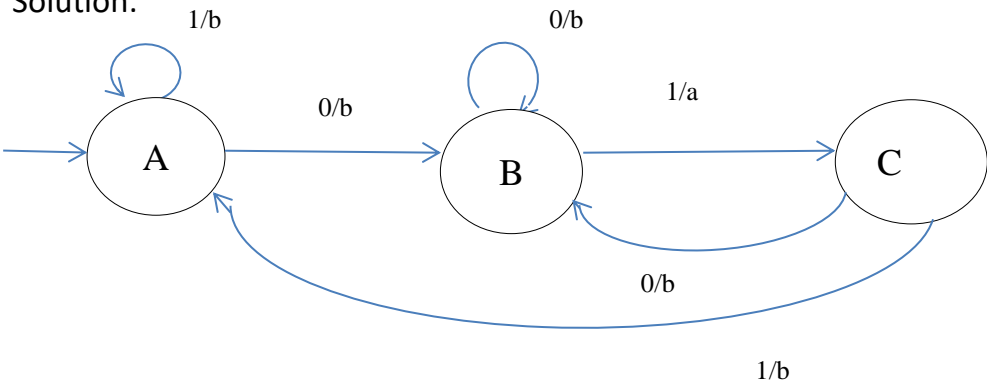
**Example2:** convert a Moore Machine to a Mealy Machine .



**Transition table for a Moore machine:**

$\Sigma$	0	1	output
Q			
A	B	A	b
B	B	C	b
C	B	A	a

**Solution:**



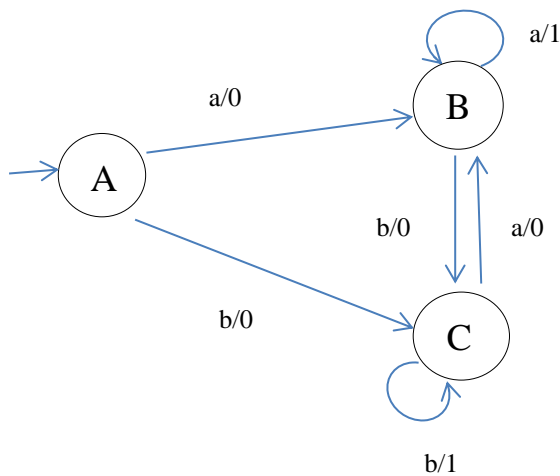
**Transition table for a Mealy machine:**

$\Sigma$ Q	<b>0</b>	<b>1</b>
<b>A</b>	B,b	A,b
<b>B</b>	B,b	C,a
<b>C</b>	B,b	A,b

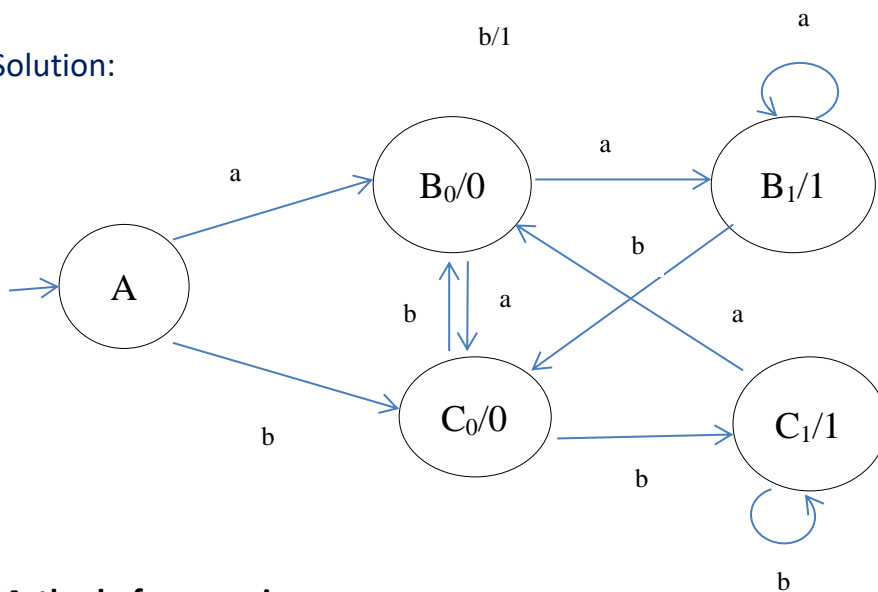
**Theorem:** let  $M1 = (Q, \Sigma, \Delta, \delta, \lambda, q_0)$  be a Mealy Machine, then there is a Moore machine  $M2$  equivalent to  $M1$ .

**Proof:** Let  $M2 = (Q \times \Delta, \Sigma, \Delta, \delta', \lambda', [q_0, b_0])$ , where  $b_0$  is an arbitrary selected member of  $\Delta$ . That is, the states of  $M2$  are pairs  $[q, b]$  consisting of a state of  $M1$  and an output symbol, in two examples below we will declare method of conversion .

**Example1:** convert the following Mealy machine to its equivalent Moore Machine.



Solution:



**Method of conversion :**

$$\lambda'([q,b],a) = [\lambda(\delta(q,a)), \lambda(q,a)]$$

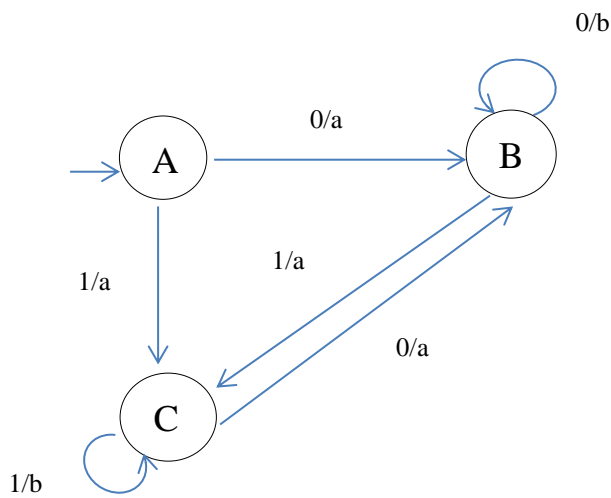
$$\lambda'([q,b]) = b$$

$$\Delta = \{\text{output symbol}\}$$

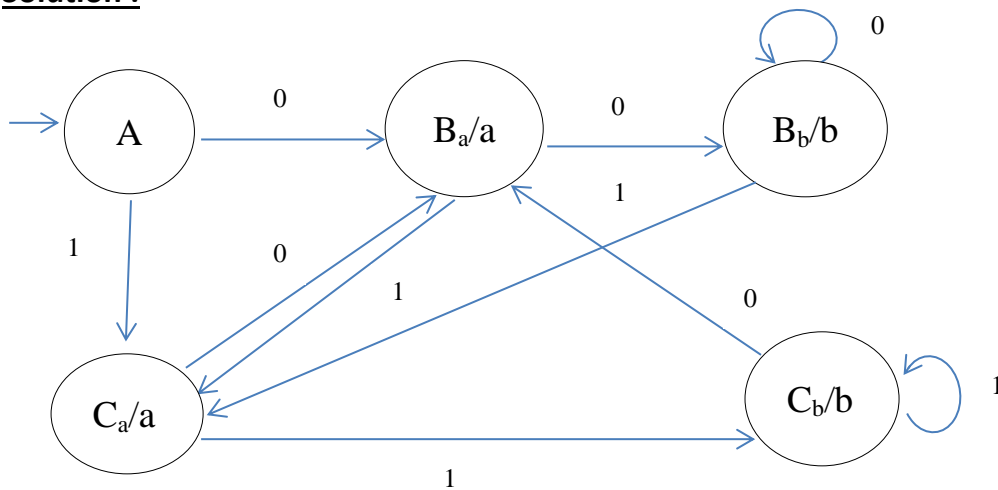
**Note :** we see every state have output except **A** does not have any output because the state **A** in original machine did not have any incoming edges.



**Example2:** convert the following Mealy machine to its equivalent Moore Machine.



**Solution :**



**Note :**

- When we convert a Moore machine to a Mealy machine → number of states were same .
- When we convert a Mealy machine to a Moore machine → number of states increased such that:

In a Mealy to Moore:  $x$  and  $y \rightarrow (x.y)$  no. of states at max

